

# COMPUTER NETWORKS

Sixth Edition

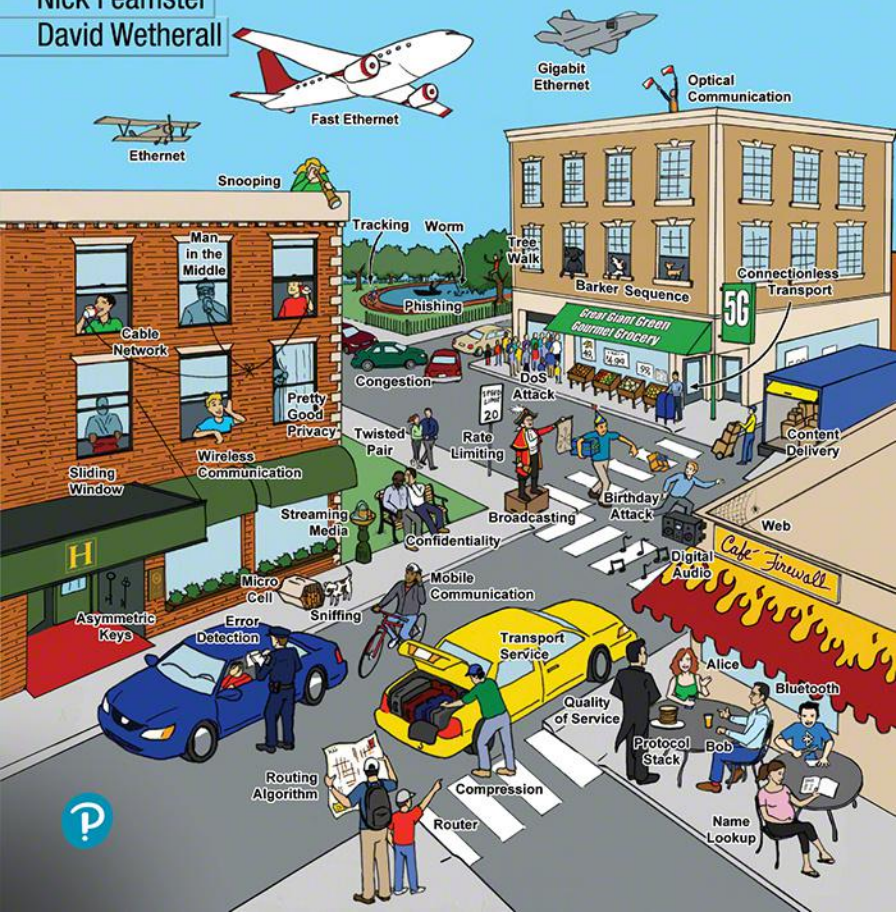
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*To Suzanne, Barbara, Daniel, Aron, Nathan, Marvin, Matilde, Olivia, and Mirte (AST)*

*To Marshini, Mila, and Kira (NF)*

*To Katrin, Lucy, and Pepper (DJW)*



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# PREFACE

This book is now in its sixth edition. Each edition has corresponded to a different phase in the way computer networks were used. When the first edition appeared in 1980, networks were an academic curiosity. When the second edition appeared in 1988, networks were used by universities and large businesses. When the third edition appeared in 1996, computer networks, especially the Internet, had become a daily reality for millions of people. By the fourth edition, in 2003, wireless networks and mobile computers had become commonplace for accessing the Web and the Internet. By the fifth edition, networks were about content distribution (especially videos using CDNs and peer-to-peer networks) and mobile phones. Now in the sixth edition, industry emphasis is on very high performance, with 5G cellular networks, 100-gigabit Ethernet, and 802.11ax WiFi at speeds up to 11 Gbps just around the corner.

## **New in the Sixth Edition**

Among the many changes in this book, the most important one is the addition of Prof. Nick Feamster as a co-author. Prof. Feamster has a Ph.D. from M.I.T. and is now a full professor at the University of Chicago.

Another important change is that Chapter 8 (on security) has been very heavily modified by Prof. Herbert Bos of the Vrije Universiteit in Amsterdam. The focus has moved from cryptography to network security. The issues of hacking, DoS attacks and so much more is front-and-center in the news almost every day, so we are very grateful that Prof. Bos has redone the chapter to deal with these important issues in detail. The chapter discusses vulnerabilities, how to fix them, how hackers respond to the fixes, how the defenders react, and so on ad infinitum. The material on cryptography has been reduced somewhat to make room for the large amount of new material on network security.

Of course, the book also has many other changes to keep up with the ever-changing world of computer networks. A chapter-by-chapter list of the major changes follows.

Chapter 1 serves the same introductory function as in previous editions, but the contents have been revised and brought up to date. Specific updates including adding additional discussions on the Internet of Things and modern cellular architectures, including 4G and 5G networks. Much of the discussion on Internet policy has also been updated, particularly the discussion on net neutrality.

Chapter 2 has been updated to include discussion of more prevalent physical media in access networks, such as DOCSIS and fiber architectures. Treatment of modern cellular network architectures and technologies was added, and the section on satellite networks was also substantially updated. Emerging technologies such as virtualization were added, including discussions on mobile virtual network operators and cellular network slicing. The policy discussion was reorganized and updated to include discussion on policy questions in the wireless arena, such as spectrum.

Chapter 3 has been updated to include DOCSIS as a protocol example, as it is a widely used access technology. Much of the error correction codes are, of course, timeless.

Chapter 4 has been brought up to date, with new material on 40- and 100-gigabit Ethernet, 802.11ac, 802.11ad, and 802.11ax. New material has been added on DOCSIS, explaining the MAC sublayer in cable networks. The material on 802.16 has been removed as it now appears that this technology is going to lose out to the cellular 4G and 5G technologies. The section on RFID has also been removed to make space for new material, but also because it was not directly network related.

Chapter 5 has been updated to clarify and modernize the discussions on congestion management. The sections on traffic management have been updated and clarified, and the discussions on traffic shaping and traffic engineering have been updated. The chapter includes an entirely new section on software-defined networking (SDN), including OpenFlow and programmable hardware (e.g., Tofino). The chapter also includes discussion on emerging applications of SDN, such as in-band network telemetry. Some of the discussion on IPv6 has also been updated.

Chapter 6 has been extensively edited to include new material on modern transport protocols, including TCP CUBIC, QUIC, and BBR. The material on performance measurement has been completely rewritten to focus on the measurement of throughput in computer networks, including an extensive discussion on the challenges of measuring access network throughput as speeds in access ISPs increase. The chapter also includes new material on measuring user quality of experience, an emerging area in performance measurement.

Chapter 7 has been heavily edited. Over 60 pages of material that is no longer relevant to a book on computer networks has been removed. The material on DNS has been almost completely rewritten to reflect modern developments in DNS, including the ongoing trends to encrypt DNS and generally improve its privacy characteristics. Emerging protocols such as DNS-over-HTTPS and other privacy-preserving techniques for DNS are discussed. The discussion of the Web has been extensively updated, to reflect the increasing deployment of encryption on the Web,

as well as extensive privacy issues (e.g., tracking) that are now pervasive on the Web. The chapter includes a completely new section on Web privacy, more extensive discussions of modern content delivery technology (e.g., content delivery networks), and an expanded discussion on peer-to-peer networks. The section on the evolution of the Internet has also been edited to reflect trends towards distributed cloud services.

Chapter 8 has been completely overhauled. In previous editions, the focus of the security chapter was almost exclusively on information security by means of cryptography. However, cryptography is only one aspect of network security and if we look at security incidents in practice, it is generally not the aspect where the problems are. To remedy this, we added new content on security principles, fundamental attack techniques, defenses, and a wide range of systems-related security issues. Moreover, we updated the existing sections by dropping some encryption techniques that are now obsolete and introducing more modern versions of protocols and standards.

Chapter 9 contains a renewed list of suggested readings and a comprehensive bibliography.

In addition, dozens of new exercises and dozens of new references have been added.

### List of Acronyms

Computer books are full of acronyms. This one is no exception. By the time you are completely finished reading this one, the following should ring a bell: AES, AMI, ARP, ARQ, ASK, BGP, BSC, CCK, CDM, CDN, CRL, DCF, DES, DIS, DMT, DMZ, DNS, EAP, ECN, EDE, EPC, FDD, FDM, FEC, FSK, GEO, GSM, HFC, HLR, HLS, HSS, IAB, IDS, IGP, IKE, IPS, ISM, ISO, ISP, ITU, IXC, IXP, KDC, LAN, LCP, LEC, LEO, LER, LLD, LSR, LTE, MAN, MEO, MFJ, MGW, MIC, MME, MPD, MSC, MSS, MTU, NAP, NAT, NAV, NCP, NFC, NIC, NID, NRZ, ONF, OSI, PAR, PCF, PCM, PCS, PGP, PHP, PIM, PKI, PON, POP, PPP, PSK, RAS, RCP, RED, RIP, RMT, RNC, RPC, RPR, RTO, RTP, SCO, SDH, SDN, SIP, SLA, SNR, SPE, SSL, TCG, TCM, TCP, TDM, TLS, TPM, UDP, URL, USB, UTP, UWB, VLR, VPN, W3C, WAF, WAN, WDM, WEP, WFQ and WPA. But don't worry. Each will appear in **boldface type** and be carefully defined before it is used. As a fun test, see how many you can identify *before* reading the book, write the number in the margin, then try again *after* reading the book.

### Instructors' Resource Materials

The following protected instructors' resource materials are available on the publisher's Web site at [www.pearsonhighered.com/tanenbaum](http://www.pearsonhighered.com/tanenbaum). For a username and password, please contact your local Pearson representative.

- Solutions manual
- PowerPoint lecture slides

**Students' Resource Materials**

Resources for students are available through the open-access Companion Web site link on [www.pearsonhighered.com/tanenbaum](http://www.pearsonhighered.com/tanenbaum), including

- Figures, tables, and programs from the book
- Steganography demo
- Protocol simulators

In addition, the authors have a Web site with other resources for students at [www.computernetworksbook.com](http://www.computernetworksbook.com).

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ANDREW S. TANENBAUM

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# 1

## INTRODUCTION

Each of the past three centuries was dominated by a single new technology. The 18th century was the era of the great mechanical systems accompanying the Industrial Revolution. The 19th century was the age of the steam engine. During the 20th century, the key technology was information gathering, processing, and distribution. Among other developments, we saw the deployment of worldwide telephone networks, the invention of radio and television, the birth and unprecedented growth of the computer industry, the launching of communication satellites, and, of course, the Internet. Who knows what miracles the 21st century will bring?

As a result of this rapid technological progress, these areas are rapidly converging in the 21st century, and the differences between collecting, transporting, storing, and processing information are quickly disappearing. Organizations with hundreds of offices spread over a wide geographical area routinely expect to be able to examine the current status of even their most remote outpost at the push of a button. As our ability to gather, process, and distribute information grows, the demand for more sophisticated information processing grows even faster.

### 1.1 USES OF COMPUTER NETWORKS

Although the computing industry is still young compared to other technical industries such as automobiles and air transportation, computers have made spectacular progress in a short time. During the first two decades of their existence,



computer systems were highly centralized, usually within a single room. Often, this room had glass windows, through which visitors could gawk at the great electronic wonder inside. A medium-sized company or university might have had one or two computers, while large institutions had at most a few dozen. The idea that within fifty years vastly more powerful computers smaller than postage stamps would be mass produced by the billions was science fiction.

The convergence of computers and communications has had a profound influence on the organization of computer systems. The once-dominant concept of the “computer center” as a room with a single large computer to which users bring their work for processing is now obsolete (although data centers holding hundreds of thousands of Internet servers are common). The old model of a single computer serving all of the organization’s computational needs has been replaced by one in which a large number of separate but interconnected computers do the job. These systems are called **computer networks**. The design and organization of these networks are the subjects of this book.

Throughout the book, we will use the term “computer network” to mean a collection of interconnected, autonomous computing devices. Two computers are said to be interconnected if they can exchange information. Interconnection can take place over a variety of transmission media including copper wire, fiber optic cable, and radio waves (e.g., microwave, infrared, communication satellites). Networks come in many sizes, shapes, and forms, as we will explore throughout the book. They are usually connected to make larger networks, with the **Internet** being the most well-known example of a network of networks.

### 1.1.1 Access to Information

Access to information comes in many forms. A common method of accessing information via the Internet is using a Web browser, which allows a user to retrieve information from various Web sites, including increasingly popular social media sites. Mobile applications on smartphones now also allow users to access remote information. Topics include the arts, business, cooking, government, health, history, hobbies, recreation, science, sports, travel, and many others. Fun comes in too many ways to mention, plus some ways that are better left unmentioned.

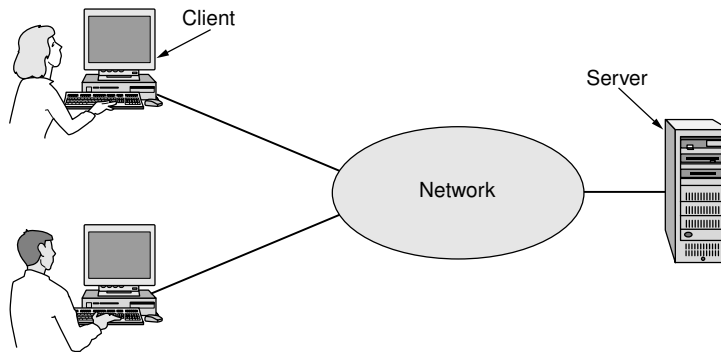
News organizations have largely migrated online, with some even ceasing print operations entirely. Access to information, including the news, is increasingly personalizable. Some online publications even allow you to tell them that you are interested in corrupt politicians, big fires, scandals involving celebrities, and epidemics, but no football, thank you. This trend certainly threatens the employment of 12-year-old paperboys, but online distribution has allowed the distribution of news to reach far larger and broader audiences.

Increasingly, news is also being curated by social media platforms, where users can post and share news content from a variety of sources, and where the news that any given user sees is prioritized and personalized based on both explicit user

preferences and complex machine learning algorithms that predict user preferences based on the user's history. Online publishing and content curation on social media platforms supports a funding model that depends largely on highly targeted behavioral advertising, which necessarily implies gathering data about the behavior of individual users. This information has sometimes been misused.

Online digital libraries and retail sites now host digital versions of content ranging from academic journals to books. Many professional organizations, such as the ACM ([www.acm.org](http://www.acm.org)) and the IEEE Computer Society ([www.computer.org](http://www.computer.org)), already have all their journals and conference proceedings online. Electronic book readers and online libraries may someday make printed books obsolete. Skeptics should take note of the effect the printing press had on the medieval illuminated manuscript.

Much information on the Internet is accessed using a client-server model, where a client explicitly requests information from a server that hosts that information, as illustrated in Fig. 1-1.

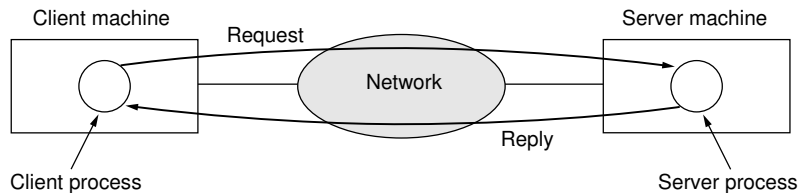


**Figure 1-1.** A network with two clients and one server.

The **client-server model** is widely used and forms the basis of much network usage. The most popular realization is that of a **Web application**, where a server generates Web pages based on its database in response to client requests that may update the database. The client-server model is applicable not only when the client and server are both in the same building (and belong to the same company), but also when they are far apart. For example, when a person at home accesses a page on the World Wide Web, the same model is employed, with the remote Web server being the server and the user's personal computer being the client. Under most conditions, one server can handle a large number (hundreds or thousands) of clients simultaneously.

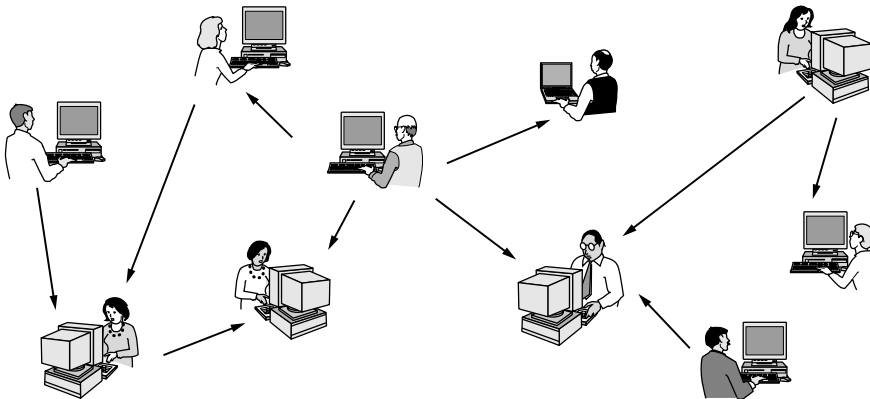
If we look at the client-server model, to a first approximation we see that two processes (running programs) are involved, one on the client machine and one on the server machine. Communication takes the form of the client process sending a

message over the network to the server process. The client process then waits for a reply message. When the server process gets the request, it performs the requested work or looks up the requested data and sends back a reply. These messages are shown in Fig. 1-2.



**Figure 1-2.** The client-server model involves requests and replies.

Another popular model for accessing information is **peer-to-peer** communication (Parameswaran et al., 2001). In this form, individuals who form a loose group can communicate with others in the group, as shown in Fig. 1-3. Every person can, in principle, communicate with one or more other people; there is no fixed division into clients and servers.



**Figure 1-3.** In a peer-to-peer system, there are no fixed clients and servers.

Many peer-to-peer systems, such as BitTorrent (Cohen, 2003), do not have a central database of content. Instead, each user maintains a local database of content, as well as a list of other members of the system. A new user can then go to any existing member to see what he has and get the names of other members to inspect for more content and more names. This lookup process can be repeated indefinitely to build up a large local database of what is out there. It is an activity that would get tedious for people, but computers excel at it.

Peer-to-peer communication is often used to share music and videos. It really hit the big time around 2000 with a music sharing service called Napster, which was shut down after a monumental copyright infringement case (Lam and Tan, 2001; and Macedonia, 2000). Legal applications for peer-to-peer communication now exist. These include fans sharing public domain music, families sharing photos and movies, and users downloading public software packages. In fact, one of the most popular Internet applications of all, email, is (conceptually) peer-to-peer. This form of communication is likely to grow considerably in the future.

### 1.1.2 Person-to-Person Communication

Person-to-person communication is the 21st century's answer to the 19th century's telephone. Email is already used on a daily basis by millions of people all over the world and its use is growing rapidly. It already routinely contains audio and video as well as text and pictures. Smell may take a while.

Many Internet users now rely on some form of **instant messaging** to communicate with other people on the Internet. This facility, derived from the UNIX *talk* program in use since around 1970, allows two people to type messages at each other in real time. There are also multi-person messaging services too, such as the **Twitter** service, which lets people send short messages (possibly including video) called “tweets” to their circle of friends or other followers or the whole world.

The Internet can be used by applications to carry audio (e.g., Internet radio stations, streaming music services) and video (e.g., Netflix, YouTube). Besides being an inexpensive way to communicate with your distant friends, these applications can provide rich experiences such as distance learning, meaning attending 8 A.M. classes without the inconvenience of having to get out of bed first. In the long run, the use of networks to enhance human-to-human communication may prove more important than any of the others. It may become hugely important to people who are geographically challenged, giving them the same access to services as people living in the middle of a big city.

Between person-to-person communications and accessing information are **social network** applications. In these applications, the flow of information is driven by the relationships that people declare between each other. One of the most popular social networking sites is **Facebook**. It lets people create and update their personal profiles and shares the updates with other people who they have declared to be their friends. Other social networking applications can make introductions via friends of friends, send news messages to friends, such as Twitter above, and much more.

Even more loosely, groups of people can work together to create content. A **wiki**, for example, is a collaborative Web site that the members of a community edit. The most famous wiki is the **Wikipedia**, an encyclopedia anyone can read or edit, but there are thousands of other wikis.

### 1.1.3 Electronic Commerce

Online shopping is already popular; users can browse the online catalogs of thousands of companies and have products shipped right to their doorsteps. After the customer buys a product electronically but cannot figure out how to use it, online technical support may be consulted.

Another area in which e-commerce is widely used is access to financial institutions. Many people already pay their bills, manage their bank accounts, and even handle their investments electronically. Financial technology or “fintech” applications allow users to conduct a wide variety of financial transactions online, including transferring money between bank accounts, or even between friends.

Online auctions of second-hand goods have become a massive industry. Unlike traditional e-commerce, which follows the client-server model, online auctions are peer-to-peer in the sense that consumers can act as both buyers and sellers, although there is a central server that holds the database of products for sale.

Some of these forms of e-commerce have acquired cute little tags based on the fact that “to” and “2” are pronounced the same. The most popular ones are listed in Fig. 1-4.

Tag	Full name	Example
B2C	Business-to-consumer	Ordering books online
B2B	Business-to-business	Car manufacturer ordering tires from a supplier
G2C	Government-to-consumer	Government distributing tax forms electronically
C2C	Consumer-to-consumer	Auctioning second-hand products online
P2P	Peer-to-peer	Music or file sharing; Skype

**Figure 1-4.** Some forms of e-commerce.

### 1.1.4 Entertainment

Our fourth category is entertainment. This has made huge strides in the home in recent years, with the distribution of music, radio and television programs, and movies over the Internet beginning to rival that of traditional mechanisms. Users can find, buy, and download MP3 songs and high-definition movies and add them to their personal collection. TV shows now reach many homes via **IPTV (IP Television)** systems that are based on IP technology instead of cable TV or radio transmissions. Media streaming applications let users tune to Internet radio stations or watch recent episodes of their favorite TV shows or movies. Naturally, all of this content can be moved around your house between different devices, displays, and speakers, usually via a wireless network.

Soon, it may be possible to search for any movie or television program ever made, in any country, and have it be displayed on your screen instantly. New films

may become interactive, where the user is occasionally prompted for the story direction (should Macbeth murder the king or just bide his time?) with alternative scenarios provided for all cases. Live television may also become interactive, with the audience participating in quiz shows, choosing among contestants, and so on.

Another form of entertainment is game playing. Already we have multi-person real-time simulation games, like hide-and-seek in a virtual dungeon, and flight simulators with the players on one team trying to shoot down the players on the opposing team. Virtual worlds provide a persistent setting in which thousands of users can experience a shared reality with three-dimensional graphics.

### 1.1.5 The Internet of Things

**Ubiquitous computing** entails computing that is embedded in everyday life, as in the vision of Mark Weiser (1991). Many homes are already wired with security systems that include door and window sensors. Also, there are many more sensors that can be folded into a smart home monitor, such as energy consumption. Smart electricity, gas, and water meters report usage over the network. This functionality saves the company money as there is then no need to send people to read the meters. Smoke detectors can call the fire department instead of just making a big noise (which has little value if no one is home). Smart refrigerators could order more milk when it is almost gone. As the cost of sensing and communication drops, more and more measurement and reporting will be done with networks. This ongoing revolution, often referred to as the **IoT (Internet of Things)**, is poised to connect just about every electronic device we purchase to the Internet.

Increasingly, consumer electronic devices are networked. For example, some high-end cameras already have a wireless network capability and use it to send photos to a nearby display for viewing. Professional sports photographers can also send their photos to their editors in real-time, first wirelessly to an access point then over the Internet. Devices such as televisions that plug into the wall can use **power-line networks** to send information throughout the house over the wires that carry electricity. It may not be very surprising to have these objects on the network, but objects that we do not think of as computers may sense and communicate information too. For example, your shower may record water usage, give you visual feedback while you lather up, and report to a home environmental monitoring application when you are done to help save on your water bill.

## 1.2 TYPES OF COMPUTER NETWORKS

There are many distinct types of computer networks. This section provides an overview of a few of these networks, including those we commonly use to access the Internet (mobile and broadband access networks); those that house the data and

applications we use every day (data-center networks); those that connect access networks to data centers (transit networks); and those that we use on a campus, office building, or other organization (enterprise networks).

### 1.2.1 Broadband Access Networks

In 1977, Ken Olsen was president of the Digital Equipment Corporation, then the number two computer vendor in the world (after IBM). When asked why Digital was not going after the personal computer market in a big way, he said: “There is no reason for any individual to have a computer in his home.” History showed otherwise and Digital no longer exists. People initially bought computers for word processing and games. Now the prevailing reason to buy a home computer is to get Internet access. Also, many consumer electronic devices, such as set-top boxes, game consoles, television sets, and even door locks, come with embedded computers that access computer networks, especially wireless networks. Home networks are broadly used for entertainment, including listening to, looking at, and creating music, photos, and videos.

Internet access provides home users with **connectivity** to remote computers. As with companies, home users can access information, communicate with other people, and buy products and services. The main benefit now comes from connecting these devices to other destinations outside of the home. Bob Metcalfe, the inventor of Ethernet, hypothesized that the value of a network is proportional to the square of the number of users because this is roughly the number of different connections that may be made (Gilder, 1993). This hypothesis is known as “Metcalfe’s law.” It helps to explain how the tremendous popularity of the Internet comes from its size.

Today, broadband access networks are proliferating. In many parts of the world, broadband access is delivered to homes through copper (e.g., telephone lines), coaxial cable (e.g., cable), or optical fiber. The speeds of broadband Internet access continue to increase as well, with many broadband access providers in developed countries delivering a gigabit per second to individual homes. In some parts of the world, particularly in developing regions, the predominant mode of Internet access is mobile.

### 1.2.2 Mobile and Wireless Access Networks

Mobile computers, such as laptops, tablets, and smartphones, are one of the fastest-growing segments of the computer industry. Their sales have already overtaken those of desktop computers. Why would anyone want one? People on the go often want to use their mobile devices to read and send email, tweet, watch movies, download music, play games, look at maps, or simply to surf the Web for information or fun. They want to do all of the things they do at home and in the office. Naturally, they want to do them from anywhere on land, sea, or in the air.

Connectivity to the Internet enables many of these mobile uses. Since having a wired connection is impossible in cars, boats, and airplanes, there is a lot of interest in wireless networks. Cellular networks operated by telephone companies are one familiar kind of wireless network that blankets us with coverage for mobile phones. Wireless **hotspots** based on the 802.11 standard are another kind of wireless network for mobile computers and portable devices such as phones and tablets. They have sprung up everywhere that people go, resulting in a patchwork of coverage at cafes, hotels, airports, schools, trains, and planes. Anyone with a mobile device and a wireless modem can just turn on their computer and be connected to the Internet through the hotspot as though the computer were plugged into a wired network.

Wireless networks are of great value to fleets of trucks, taxis, delivery vehicles, and repair-persons for keeping in contact with their home base. For example, in many cities, taxi drivers are independent businessmen, rather than being employees of a taxi company. In some of these cities, the taxis have a display the driver can see. When a customer calls up, a central dispatcher types in the pickup and destination points. This information is displayed on the drivers' displays and a beep sounds. The first driver to hit a button on the display gets the call. The rise of mobile and wireless networking has also led to a revolution in ground transportation itself, with the "sharing economy" allowing drivers to use their on phones as a dispatch device, as with ride-sharing companies such as Uber and Lyft.

Wireless networks are also important to the military. If you have to be able to fight a war anywhere on Earth at short notice, counting on using the local networking infrastructure is probably not a good idea. It is better to bring your own.

Although wireless networking and mobile computing are often related, they are not identical, as Fig. 1-5 shows. Here, we see a distinction between **fixed wireless** and **mobile wireless** networks. Even notebook computers are sometimes wired. For example, if a traveler plugs a laptop computer into the wired network jack in a hotel room, he has mobility without a wireless network. The growing pervasiveness of wireless networks is making this situation increasingly rare, although for high performance, wired networks are always better.

Wireless	Mobile	Typical applications
No	No	Desktop computers in offices
No	Yes	A laptop computer used in a hotel room
Yes	No	Networks in unwired buildings
Yes	Yes	Store inventory with a handheld computer

**Figure 1-5.** Combinations of wireless networks and mobile computing.

Conversely, some wireless computers are not mobile. In people's homes, and in offices or hotels that lack suitable cabling, it can be more convenient to connect desktop computers or media players wirelessly than to install wires. Installing a



wireless network may require simply buying a small box with some electronics in it, unpacking it, and plugging it in. This solution may be far cheaper than having workmen put in cable ducts to wire the building.

Finally, there are also true mobile, wireless applications, such as people walking around stores with handheld computers recording inventory. At many busy airports, car rental return clerks work in the parking lot with wireless mobile computers. They scan the barcodes or RFID chips of returning cars, and their mobile device, which has a built-in printer, calls the main computer, gets the rental information, and prints out the bill on the spot.

A key driver of mobile, wireless applications is the mobile phone. The convergence between telephones and the Internet is accelerating the growth of mobile applications. **Smartphones**, such as Apple's iPhone and Samsung's Galaxy, combine aspects of mobile phones and mobile computers. These phones connect to wireless hotspots, too, and automatically switch between networks to choose the best option for the user. **Text messaging** or **texting** (or **Short Message Service** as it is known outside the U.S.) over the cellular network was tremendously popular at its outset. It lets a mobile phone user type a short message that is then delivered by the cellular network to another mobile subscriber. Texting is extremely profitable since it costs the carrier but a tiny fraction of one cent to relay a text message, a service for which it charges far more. Typing short text messages on mobile phones was, for a time, an immense money maker for mobile carriers. Now, many alternatives that use either the phone's cellular data plan or wireless network, including WhatsApp, Signal, and Facebook Messenger, have overtaken SMS.

Other consumer electronics devices can also use cellular and hotspot networks to stay connected to remote computers. Tablets and electronic book readers can download a newly purchased book or the next edition of a magazine or today's newspaper wherever they roam. Electronic picture frames can update their displays on cue with fresh images.

Mobile phones typically know their own locations. **GPS (Global Positioning System)** can directly locate a device, and mobile phones often also triangulate between Wi-Fi hotspots with known locations to determine their location. Some applications are location-dependent. Mobile maps and directions are an obvious candidate as your GPS-enabled phone and car probably have a better idea of where you are than you do. So, too, are searches for a nearby bookstore or Chinese restaurant, or a local weather forecast. Other services may record location, such as annotating photos and videos with the place at which they were made. This annotation is known as **geo-tagging**.

Mobile phones are being increasingly used in **m-commerce (mobile-commerce)** (Senn, 2000). Short text messages from the mobile are used to authorize payments for food in vending machines, movie tickets, and other small items instead of cash and credit cards. The charge then appears on the mobile phone bill. When equipped with **NFC (Near Field Communication)**, technology the mobile can act as an RFID smartcard and interact with a nearby reader for payment. The

driving forces behind this phenomenon are the mobile device makers and network operators, who are trying hard to figure out how to get a piece of the e-commerce pie. From the store's point of view, this scheme may save them most of the credit card company's fee, which can be several percent. Of course, this plan may backfire, since customers in a store might use the RFID or barcode readers on their mobile devices to check out competitors' prices before buying and use them to get a detailed report on where else an item can be purchased nearby and at what price.

One huge thing that m-commerce has going for it is that mobile phone users are accustomed to paying for everything (in contrast to Internet users, who expect everything to be free). If an Internet Web site charged a fee to allow its customers to pay by credit card, there would be an immense bellowing from the users. If, however, a mobile phone operator let its customers pay for items in a store by waving the phone at the cash register and then tacks on a small fee for this convenience, it would probably be accepted as normal. Time will tell.

The uses of mobile and wireless computers will grow rapidly in the future as the size of computers shrinks, probably in ways no one can now foresee. Let us take a quick look at some possibilities. **Sensor networks** have nodes that gather and relay information they sense about the state of the physical world. The nodes may be embedded in familiar devices such as cars or phones, or they may be small separate devices. For example, your car might gather data on its location, speed, vibration, and fuel efficiency from its on-board diagnostic system and upload this information to a database (Hull et al., 2006). Those data can help find potholes, plan trips around congested roads, and tell you if you are a "gas guzzler" compared to other drivers on the same stretch of road.

Sensor networks are revolutionizing science by providing a wealth of data on behavior that could not previously be observed. One example is tracking the migration of individual zebras by placing a small sensor on each animal (Juang et al., 2002). Researchers have packed a wireless computer into a single square cubic millimeter (Warneke et al., 2001). With mobile computers this small, even small birds, rodents, and insects can be tracked.

Wireless parking meters can accept credit or debit card payments with instant verification over the wireless link. They can also report when they are in use, which can let drivers download a recent parking map to their car so they can find an available spot more easily. Of course, when a meter expires, it might also check for the presence of a car (by bouncing a signal off it) and report the expiration to parking enforcement. It has been estimated that city governments in the U.S. alone could collect an additional \$10 billion this way (Harte et al., 2000).

### 1.2.3 Content Provider Networks

Many Internet services are now served from "the cloud," or a **data-center network**. Modern data center networks have hundreds of thousands or millions of servers in a single location, usually in a very dense configuration of rows of racks

in buildings that can be more than a kilometer long. Data center networks serve the increasingly growing demands of **cloud computing** and are designed to move large amounts of data between servers in the data center, as well as between the data center and the rest of the Internet.

Today, many of the applications and services you use, ranging from the Web sites you visit to the cloud-based document editor you use to take notes, store data in a data center network. Data center networks face challenges of scale, both for network throughput and for energy usage. One of the main network throughput challenges is the so-called “cross-section bandwidth,” which is the data rate that can be delivered between any two servers in the network. Early data-center network designs were based on a simple tree topology, with three layers of switches: access, aggregate, and core; this simple design did not scale well, and was also to be subject to faults.

Many popular Internet services need to deliver content to users around the world. To do so, many sites and services on the Internet use a **CDN (Content Delivery Network)**. A CDN is a large collection of servers that are geographically distributed in such a way that content is placed as close as possible to the users that are requesting it. Large content providers such as Google, Facebook, and Netflix operate their own CDNs. Some CDNs, such as Akamai and Cloudflare, offer hosting services to smaller services that do not have their own CDN.

Content that users want to access, ranging from static files to streaming video, may be replicated in many locations across a single CDN. When a user requests content, the CDN must decide which replica it should serve to that user. This process must consider the distance from each replica to the client, the load on each CDN server, and traffic load and congestion on the network itself.

### 1.2.4 Transit Networks

Internet travels over many independently operated networks. The network run by your Internet service provider is typically not the same network as the one that hosts the content for the Web sites that you commonly visit. Typically, content and applications are hosted in data-center networks, and you may be accessing that content from an access network. Content must thus traverse the Internet from the data center to the access network, and ultimately to your device.

When the content provider and your **ISP (Internet Service Provider)** are not directly connected, they often rely on a **transit network** to carry the traffic between them. Transit networks typically charge both the ISP and the content provider for carrying traffic from end-to-end. If the network hosting the content and the access network exchange enough traffic between them, they may decide to interconnect directly. One example where direct interconnection is common is between large ISPs and large content providers, such as Google or Netflix. In these cases, the ISP and the content provider must build and maintain network infrastructure to facilitate interconnecting directly, often in many geographic locations.

Transit networks are traditionally called **backbone networks** because they have had the role of carrying traffic between two endpoints. Many years ago, transit networks were hugely profitable because every other network would rely on them (and pay them) to connect to the rest of the Internet.

The last decade, however, has witnessed two trends. The first trend is the consolidation of content in a handful of large content providers, spawned by the proliferation of cloud-hosted services and large content delivery networks. The second trend is the expansion of the footprint of individual access ISP networks: whereas access ISPs may have once been small and regional, many access ISPs have national (or even international) footprints, which has increased both the range of geographic locations where they can connect to other networks as well as their subscriber base. As the size (and negotiating power) of the access networks and the content provider networks continues to increase, the larger networks have come to rely less on transit networks to deliver their traffic, preferring often to directly interconnect and rely on the transit network only as a backup.

### 1.2.5 Enterprise Networks

Most organizations (e.g., companies, universities) have many computers. Each employee may use a computer to perform tasks ranging from product design to payroll. In the common case, these machines are connected on a common network, which allows the employees to share data, information, and compute resources with one another.

**Resource sharing** makes programs, equipment, and especially data available to other users on the network without regard to the physical location of the resource or the user. One widespread example is having a group of office workers share a common printer. Many employees do not need a private printer and a high-volume networked printer is often less expensive, faster, and easier to maintain than a large collection of individual printers.

Probably, even more important than sharing physical resources such as printers and backup systems is sharing information. Most companies have customer records, product information, inventories, financial statements, tax information, and much more online. If all of its computers suddenly went down, a bank could not last more than five minutes. A modern manufacturing plant, with a computer-controlled assembly line, would not last even five seconds. Even a small travel agency or three-person law firm is now highly dependent on computer networks for allowing employees to access relevant information and documents instantly.

For smaller companies, the computers may be located in a single office even a single building; in the case of larger companies, the computers and employees may be scattered over dozens of offices and plants in many countries. Nevertheless, a salesperson in New York might sometimes need access to a product inventory database in Singapore. Networks called **VPNs (Virtual Private Networks)** connect

the individual networks at different sites into one logical network. In other words, the mere fact that a user happens to be 15,000 km away from his data should not prevent him from using the data as though they were local. This goal may be summarized by saying that it is an attempt to end the “tyranny of geography.”

In the simplest of terms, one can imagine a company’s information system as consisting of one or more databases with company information and some number of employees who need to access them remotely. In this model, the data are stored on powerful computers called **servers**. Often, these are centrally housed and maintained by a system administrator. In contrast, the employees have simpler machines, called **clients**, on their desks, with which they access remote data, for example, to include in spreadsheets they are constructing. (Sometimes we will refer to the human user of the client machine as the “client,” but it should be clear from the context whether we mean the computer or its user.) The client and server machines are connected by a network, as illustrated in Fig. 1-1. Note that we have shown the network as a simple oval, without any detail. We will use this form when we mean a network in the most abstract sense. When more detail is required, it will be provided.

A second goal of setting up an enterprise computer network has to do with people rather than information or even computers. A computer network can provide a powerful **communication medium** among employees. Virtually every company that has two or more computers now has **email (electronic mail)**, which employees generally use for a great deal of daily communication. In fact, a common gripe around the water cooler is how much email everyone has to deal with, much of it quite meaningless because bosses have discovered that they can send the same (often content-free) message to all their subordinates at the push of a button.

Telephone calls between employees may be carried by the computer network instead of by the phone company. This technology is called **IP telephony** or **VoIP (Voice over IP)** when Internet technology is used. The microphone and speaker at each end may belong to a VoIP-enabled phone or the employee’s computer. Companies find this a wonderful way to save on their telephone bills.

Other, much richer forms of communication are made possible by computer networks. Video can be added to audio so that multiple employees at distant locations can see and hear each other as they hold a meeting. This technique is a powerful tool for eliminating the cost and time previously devoted to travel. **Desktop sharing** lets remote workers see and interact with a graphical computer screen. This makes it easy for two or more people who work far apart to read and write a shared blackboard or write a report together. When one worker makes a change to an online document, the others can see the change immediately, instead of waiting several days for a letter. Such a speedup makes cooperation among far-flung groups of people easy where it previously had been impossible. More ambitious forms of remote coordination such as telemedicine are only now starting to be used (e.g., remote patient monitoring) but may become much more important. It is

sometimes said that communication and transportation are having a race, and whichever wins will make the other obsolete.

A third goal for many companies is doing business electronically, especially with customers and also suppliers. Airlines, bookstores, and other retailers have discovered that many customers like the convenience of shopping from home. Consequently, many companies provide catalogs of their goods and services online and take orders online. Manufacturers of automobiles, aircraft, and computers, among others, buy subsystems from many suppliers and then assemble the parts. Using computer networks, manufacturers can place orders electronically as needed. This reduces the need for large inventories and enhances efficiency.

## 1.3 NETWORK TECHNOLOGY, FROM LOCAL TO GLOBAL

Networks can range from small and personal to large and global. In this section, we explore the various networking technologies that implement networks at different sizes and scales.

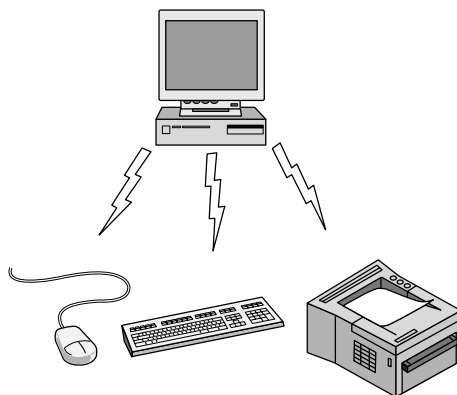
### 1.3.1 Personal Area Networks

**PANs (Personal Area Networks)** let devices communicate over the range of a person. A common example is a wireless network that connects a computer with its peripherals. Other examples include the network that connects your wireless headphones and your watch to your smartphone. It is also often used to connect a headset to a mobile phone without cords, and it can allow your digital music player to connect to your car merely being brought within range.

Almost every computer has an attached monitor, keyboard, mouse, and printer. Without using wireless, this connection must be done with cables. Many new users have so much trouble finding the right cables and plugging them into the right little holes (even though they are usually shape and color coded) that most computer vendors offer the option of sending a technician to the user's home to do it. To help these users, some companies got together to design a short-range wireless network called **Bluetooth** to connect these components without wires. The idea is that if your devices have Bluetooth, then you do not need to deal with cables. You just put them down, turn them on, and they begin communicating. For many people, this ease of operation is a big plus.

In the simplest form, Bluetooth networks use the master-slave paradigm shown in Fig. 1-6. The system unit (the PC) is normally the master, talking to the mouse or keyboard as slaves. The master tells the slaves what addresses to use, when they can transmit, how long they can transmit, what frequencies they can use, and so on. We will discuss Bluetooth in more detail in Chap. 4.

PANs can also be built with a variety of other technologies that communicate over short ranges, as we will discuss in Chap. 4.



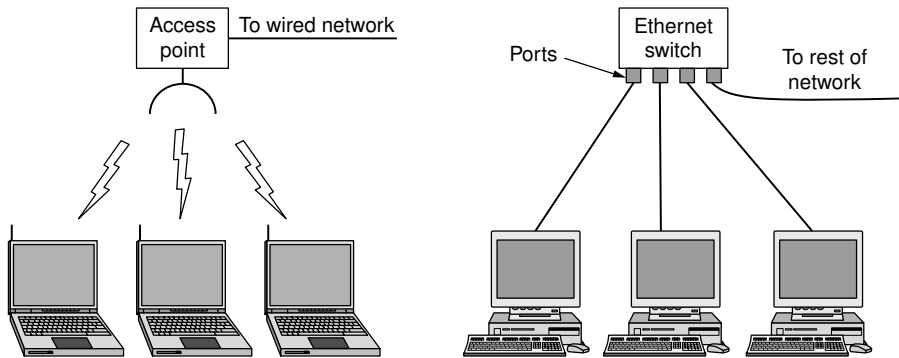
**Figure 1-6.** Bluetooth PAN configuration.

### 1.3.2 Local Area Networks

A **LAN (Local Area Network)** is a private network that operates within and nearby a single building such as a home, office, or factory. LANs are widely used to connect personal computers and consumer electronics to let them share resources (e.g., printers) and exchange information.

Wireless LANs are pervasive today. They initially gained popularity in homes, older office buildings, cafeterias, and other places where installing cables introduced too much cost. In these systems, every computer has a radio modem and an antenna that it uses to communicate with other computers. In most cases, each computer talks to a device called an **AP (Access Point)**, **wireless router**, or **base station**, as shown in Fig. 1-7(a). This device relays packets between the wireless computers and also between them and the Internet. Being the AP is like being the popular kid at school because everyone wants to talk to you. Another common scenario entails nearby devices relaying packets for one another in a so-called **mesh network** configuration. In some cases, the relays are the same nodes as the endpoints; more commonly, however, a mesh network will include a separate collection of nodes whose sole responsibility is relaying traffic. Mesh network settings are common in developing regions where deploying connectivity across a region may be cumbersome or costly. They are also becoming increasingly popular for home networks, particularly in large homes.

There is a popular standard for wireless LANs called **IEEE 802.11**, commonly called **WiFi**. It runs at speeds from 11 Mbps (802.11b) to 7 Gbps (802.11ad). Please note that in this book we will adhere to tradition and measure line speeds in megabits/sec, where 1 Mbps is 1,000,000 bits/sec, and gigabits/sec, where 1 Gbps is 1,000,000,000 bits/sec. Powers of two are used only for storage, where a 1 MB memory is  $2^{20}$  or 1,048,576 bytes. We will discuss 802.11 in Chap. 4.



**Figure 1-7.** Wireless and wired LANs. (a) 802.11. (b) Switched Ethernet.

Wired LANs use many different transmission technologies; common physical modes of transmission are copper, coaxial cable, and optical fiber. LANs have limited size, which means that the worst-case transmission time is bounded and known in advance. Knowing these bounds helps with the task of designing network protocols. Typically, wired LANs can run at speeds ranging from 100 Mbps to 40 Gbps. They also have low latency (never more than tens of milliseconds, and often much less) and transmission errors are infrequent. Wired LANs typically have lower latency, lower packet loss, and higher throughput than wireless LANs, but over time this performance gap has narrowed. It is far easier to send signals over a wire or through a fiber than through the air.

Many wired LANs comprise point-to-point wired links. IEEE 802.3, popularly called **Ethernet**, is by far the most common type of wired LAN. Fig. 1-7(b) shows an example **switched Ethernet** topology. Each computer speaks the Ethernet protocol and connects to a device called a **switch** with a point-to-point link. The job of the switch is to relay packets between computers that are attached to it, using the address in each packet to determine which computer to send it to.

A switch has multiple **ports**, each of which can connect to one other device, such as a computer or even another switch. To build larger LANs, switches can be plugged into each other using their ports. What happens if you plug them together in a loop? Will the network still work? Luckily, someone thought of this case, and now all switches in the world use her anti-looping algorithm (Perlman, 1985). It is the job of the protocol to sort out what paths packets should travel to safely reach the intended computer. We will see how this works in Chap. 4.

It is also possible to divide one large physical LAN into two smaller logical LANs. You might wonder why this would be useful. Sometimes, the layout of the network equipment does not match the organization's structure. For example, the engineering and finance departments of a company might have computers on the same physical LAN because they are in the same wing of the building, but it might be easier to manage the system if engineering and finance logically each had its



own network **VLAN (Virtual LAN)**. In this design, each port is tagged with a “color,” say green for engineering and red for finance. The switch then forwards packets so that computers attached to the green ports are separated from the computers attached to the red ports. Broadcast packets sent on a red port, for example, will not be received on a green port, just as though there were two separate physical LANs. We will cover VLANs at the end of Chap. 4.

There are other wired LAN topologies, too. In fact, switched Ethernet is a modern version of the original Ethernet design that broadcasts all packets over a single linear cable. At most one machine could successfully transmit at a time, and a distributed arbitration mechanism was used to resolve conflicts. It used a simple algorithm: computers could transmit whenever the cable was idle. If two or more packets collided, each computer just waited a random time and tried later. We will call that version **classic Ethernet** for clarity, and as you no doubt suspected, you will learn about it in Chap. 4.

Both wireless and wired broadcast LANs can allocate resources statically or dynamically. A typical static allocation would be to divide time into discrete intervals and use a round-robin algorithm, allowing each machine to broadcast only when its time slot comes up. Static allocation wastes channel capacity when a machine has nothing to transmit or receive during its allocated slot, so most systems attempt to allocate the channel dynamically (i.e., on demand).

Dynamic allocation methods for a common channel are either centralized or decentralized. In a centralized channel allocation method, there is a single entity, for example, the base station in cellular networks, which determines who goes next. It might do so by accepting multiple packets and prioritizing them according to some internal algorithm. In a decentralized channel allocation method, there is no central entity; each machine must decide for itself whether to transmit. You might think that this approach would lead to chaos, but later we will study many algorithms designed to bring order out of the potential chaos—provided, of course, that all the machines obey the rules.

### 1.3.3 Home Networks

It is worth giving specific attention to LANs in the home, or **home networks**. Home networks are a type of LAN; they may have a broad, diverse range of Internet-connected devices, and must be particularly easy to manage, dependable, and secure, especially in the hands of nontechnical users.

Many years ago, a home network would probably have consisted of a few laptops on a wireless LAN. Today, a home network may include devices such as smartphones, wireless printers, thermostats, burglar alarms, smoke detectors, lightbulbs, cameras, televisions, stereos, smart speakers, refrigerators, and so on. The proliferation of Internet-connected appliances and consumer electronics, often called the Internet of things, makes it possible to connect just about any electronic

device (including sensors of many types) to the Internet. This huge scale and diversity of Internet connected devices introduces new challenges for designing, managing, and securing a home network. Remote monitoring of the home is becoming increasingly common, with applications ranging from security monitoring to maintenance to aging in place, as many grown children are willing to spend some money to help their aging parents live safely in their own homes.

Although the home network is just another LAN, in practice it is likely to have different properties than other LANs, for several reasons. First, the devices that people connect to their home network need to be easy to install and maintain. Wireless routers were at one point very commonly returned to stores because people bought them expecting to have a wireless network work “out of the box” but instead found themselves confronted with the prospect of many calls to technical support. The devices need to be foolproof and work without requiring the user to read and fully understand a 50-page manual.

Second, security and reliability have higher stakes because insecurity of the devices may introduce direct threats to consumer health and safety. Losing a few files to an email virus is one thing; having a burglar disarm your security system from his phone and then plunder your house is something quite different. The past few years have seen countless examples of insecure or malfunctioning IoT devices that have resulted in everything from frozen pipes to remote control of devices through malicious third-party scripts. The lack of serious security on many of these devices has made it possible for an eavesdropper to observe details about user activity in the home; even when the contents of the communication are encrypted, simply knowing the type of device that is communicating and the volumes and times of traffic can reveal a lot about private user behavior.

Third, home networks evolve organically, as people buy various consumer electronics devices and connect them to the network. As a result, in contrast to a more homogeneous enterprise LAN, the set of technologies connected to the home network may be significantly more diverse. Yet, despite this diversity, people expect these devices to be able to interact (e.g., they want to be able to use the voice assistant manufactured by one vendor to control the lights from another vendor). Once installed, the devices may remain connected for years (or decades). This means no interface wars: Telling consumers to buy peripherals with IEEE 1394 (FireWire) interfaces and a few years later retracting that and saying USB 3.0 is the interface-of-the-month and then switching that to 802.11g—oops, no, make that 802.11n—no wait, 802.11ac—sorry, we mean 802.11ax, is not tenable.

Finally, profit margins are small in consumer electronics, so many devices aim to be as inexpensive as possible. When confronted with a choice about which Internet-connected digital photo frame to buy, many users may opt for the less-expensive one. The pressure to reduce consumer device costs makes achieving the above goals even more difficult. Security, reliability, and interoperability all ultimately cost money. In some cases, manufacturers or consumers may need powerful incentives to make and stick to recognized standards.

Home networks typically operate over wireless networks. Convenience and cost favors wireless networking because there are no wires to fit, or worse, retrofit. As Internet-connected devices proliferate, it becomes increasingly inconvenient to drop a wired network port everywhere in the home where there is a power outlet. Wireless networks are more convenient and more cost-effective. Reliance on wireless networks in the home, however, does introduce unique performance and security challenges. First, as users exchange more traffic on their home networks and connect more devices to them, the home wireless network is increasingly becoming a performance bottleneck. When the home network is performing poorly, a common pastime is to blame the ISP for the poor performance. ISPs tend not to like this so much.

Second, wireless radio waves can travel through walls (in the popular 2.4 GHz band, but less so at 5 GHz). Although wireless security has improved substantially over the last decade, it still has been subject to many attacks that allow eavesdropping, and certain aspects of the traffic, such as device hardware addresses and traffic volume, remain unencrypted. In Chap. 8, we will study how encryption can be used to provide security, but it is easier said than done with inexperienced users.

**Power-line networks** can also let devices that plug into outlets broadcast information throughout the house. You have to plug in the TV anyway, and this way it can get Internet connectivity at the same time. These networks carry both power and data signals at the same time; part of the solution is to run these two functions on different frequency bands.

### 1.3.4 Metropolitan Area Networks

A **MAN (Metropolitan Area Network)** covers a city. The best-known examples of MANs are the cable television networks. These systems grew from earlier community antenna systems used in areas with poor over-the-air television reception. In those early systems, a large antenna was placed on top of a nearby hill and a signal was then piped to the subscribers' houses.

At first, these networks were locally designed, ad hoc systems. Then, companies began jumping into the business, getting contracts from local governments to wire up entire cities. The next step was television programming and even entire channels designed for cable only. Often, these channels were highly specialized, such as all news, all sports, all cooking, all gardening, and so on. But from their inception until the late 1990s, they were intended for television reception only.

When the Internet began attracting a mass audience, the cable TV network operators began to realize that with some changes to the system, they could provide two-way Internet service in unused parts of the spectrum. At that point, the cable TV system began to morph from simply a way to distribute television to a metropolitan area network. To a first approximation, a MAN might look something like the system shown in Fig. 1-8. In this figure, we see both television signals and Internet being fed into the centralized **cable head-end**, (or cable modem termination

system) for subsequent distribution to people's homes. We will come back to this subject in detail in Chap. 2.

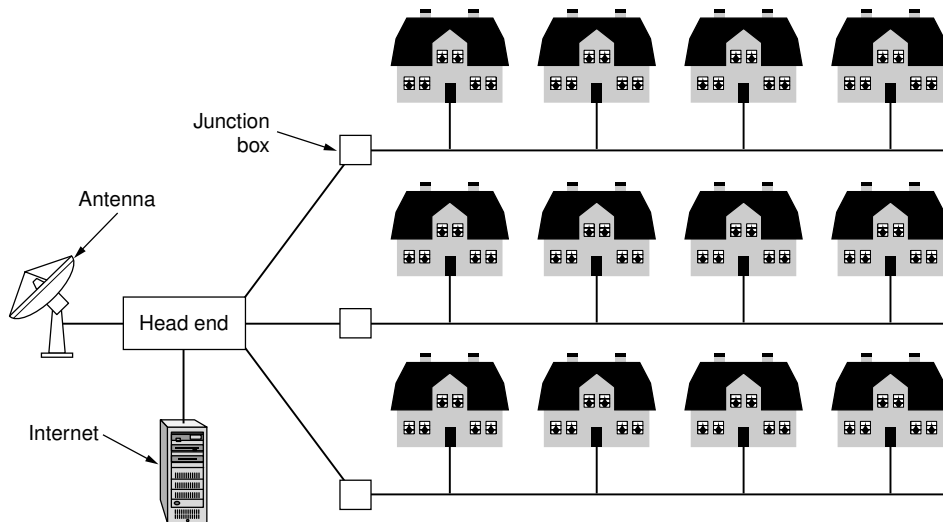


Figure 1-8. A metropolitan area network based on cable TV.

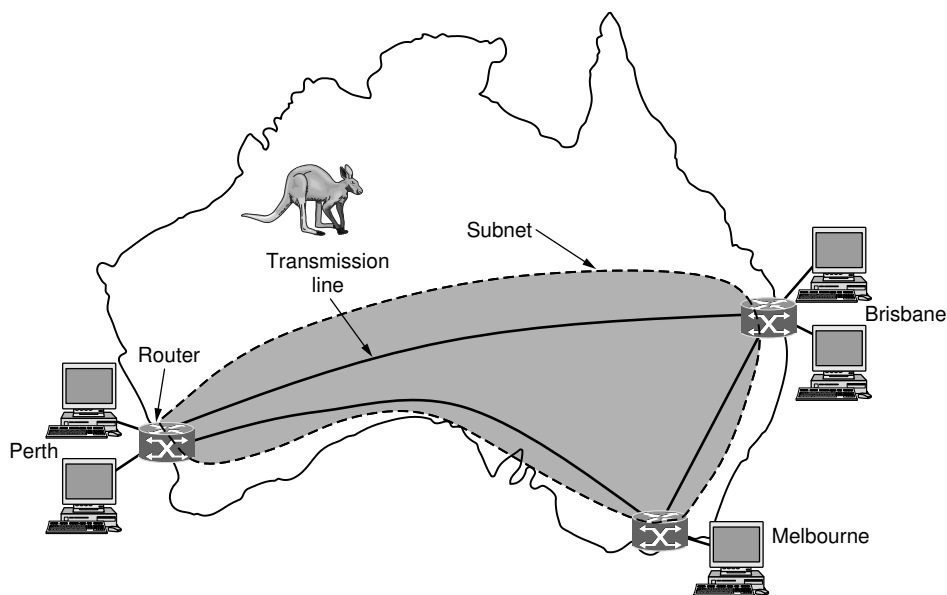
Cable television is not the only MAN. Recent developments in high-speed wireless Internet access have resulted in another MAN, which has been standardized as IEEE 802.16 and is popularly known as **WiMAX**. It does not seem to be catching on, however. Other wireless technologies, **LTE (Long Term Evolution)** and **5G**, will also be covered there.

### 1.3.5 Wide Area Networks

A **WAN (Wide Area Network)** spans a large geographical area, often a country, a continent, or even multiple continents. A WAN may serve a private organization, as in the case of an enterprise WAN, or it may be a commercial service offering, as in the case of a transit network.

We will begin our discussion with wired WANs, using the example of a company with branch offices in different cities. The WAN in Fig. 1-9 connects offices in Perth, Melbourne, and Brisbane. Each of these offices contains computers intended for running user (i.e., application) programs. We will follow conventional usage and call these machines **hosts**. The rest of the network that connects these hosts is then called the **communication subnet**, or just **subnet** for short. The subnet carries messages from host to host, just as the telephone system carries words (really just sounds) from speaker to listener.

In most WANs, the subnet consists of two distinct components: transmission lines and switching elements. **Transmission lines** move bits between machines.



**Figure 1-9.** WAN that connects three branch offices in Australia.

They can be made of copper wire, coaxial cable, optical fiber, or radio links. Most organizations do not have transmission lines lying about, so instead they use the lines from a telecommunications company. **Switching elements**, or **switches**, are specialized devices that connect two or more transmission lines. When data arrive on an incoming line, the switching element must choose an outgoing line on which to forward them. These switching computers have been called by various names in the past; the name **router** is now most commonly used. Unfortunately, some people pronounce it “router” while others have it rhyme with “doubter.” Determining the correct pronunciation will be left as an exercise for the reader. (Note: the perceived correct answer may depend on where you live.)

In most WANs, the network contains many transmission lines, each connecting a pair of routers. Two routers that do not share a transmission line must do so via other routers. There may be many paths in the network that connect these two routers. How the network makes the decision as to which path to use is called a **routing algorithm**. How each router makes the decision as to where to send a packet next is called a **forwarding algorithm**. We will study some of both types in detail in Chap. 5.

A short comment about the term “subnet” is in order here. Originally, its *only* meaning was the collection of routers and communication lines that moved packets from the source host to the destination host. Readers should be aware that it has acquired a second, more recent meaning in conjunction with network addressing.

We will discuss that meaning in Chap. 5 and stick with the original meaning (a collection of lines and routers) until then.

The WAN as we have described it looks similar to a large wired LAN, but there are some important differences that go beyond long wires. Usually in a WAN, the hosts and subnet are owned and operated by different people. In our example, the employees might be responsible for their own computers, while the company's IT department is in charge of the rest of the network. We will see clearer boundaries in the coming examples, in which the network provider or telephone company operates the subnet. Separation of the pure communication aspects of the network (the subnet) from the application aspects (the hosts) greatly simplifies the overall network design.

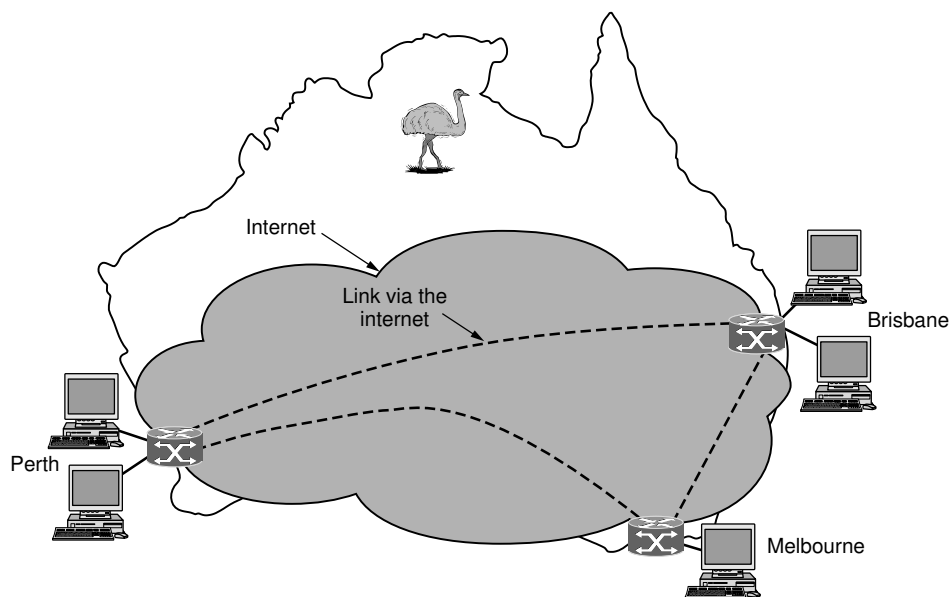
A second difference is that the routers will usually connect different kinds of networking technology. The networks inside the offices may be switched Ethernet, for example, while the long-distance transmission lines may be SONET links (which we will cover in Chap. 2). Some device needs to join them. The astute reader will notice that this goes beyond our definition of a network. This means that many WANs will in fact be **internetworks**, or composite networks that comprise more than one network. We will have more to say about internetworks in the next section.

A final difference is in what is connected to the subnet. This could be individual computers, as was the case for connecting to LANs, or it could be entire LANs. This is how larger networks are built from smaller ones. As far as the subnet is concerned, it does the same job.

### Virtual Private Networks and SD-WANs

Rather than lease dedicated transmission lines, an organization might rely on Internet connectivity to connect its offices. This allows connections to be made between the offices as virtual links that use the underlying capacity of the Internet. As mentioned earlier, this arrangement, shown in Fig. 1-10, is called a virtual private network. In contrast to a network with dedicated physical links, a VPN has the usual advantage of virtualization, which is that it provides flexible reuse of a resource (Internet connectivity). A VPN also has the usual disadvantage of virtualization, which is a lack of control over the underlying resources. With a dedicated line, the capacity is clear. With a VPN, performance may vary with that of the underlying Internet connectivity. The network itself may also be operated by a commercial Internet service provider (ISP). Fig. 1-11 shows this structure, which connects the WAN sites to each other, as well as to the rest of the Internet.

Other kinds of WANs make heavy use of wireless technologies. In satellite systems, each computer on the ground has an antenna through which it can exchange data with a satellite in orbit. All computers can hear the output *from* the satellite, and in some cases, they can also hear the upward transmissions of their



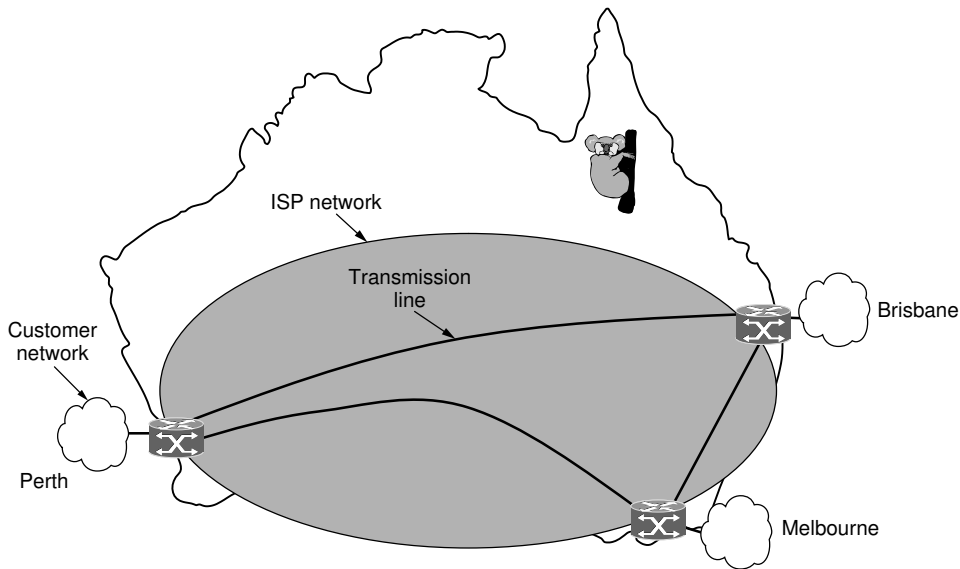
**Figure 1-10.** WAN using a virtual private network.

fellow computers *to* the satellite as well. Satellite networks are inherently broadcast and are most useful when broadcast is important or no ground-based infrastructure is present (think: oil companies exploring in an isolated desert).

The cellular telephone network is another example of a WAN that uses wireless technology. This system has already gone through five generations. The first generation was analog and for voice only. The second generation was digital and for voice only. The third generation is digital and is for both voice and data. The fourth generation is purely digital, even for voice. The fifth generation is also pure digital and much faster than the fourth, with lower delays as well.

Each cellular base station covers a distance much larger than a wireless LAN, with a range measured in kilometers rather than tens of meters. The base stations are connected to each other by a backbone network that is usually wired. The data rates of cellular networks are often on the order of 100 Mbps, much smaller than a wireless LAN that can range up to on the order of 7 Gbps. We will have a lot to say about these networks in Chap. 2.

More recently, organizations that are distributed across geographic regions and need to connect sites are designing and deploying so-called **software-defined WANs** or **SD-WANs**, which use different, complementary technologies to connect disjoint sites but provide a single **SLA (Service-Level Agreement)** across the network. For example, a network might possibly use a combination of more-expensive dedicated leased lines to connect multiple remote locations and complementary,



**Figure 1-11.** WAN using an ISP network.

less-expensive commodity Internet connectivity to connect these locations. Logic written in software reprograms the switching elements in real time to optimize the network for both cost and performance. SD-WANs are one example of an **SDN (Software-Defined Network)**, a technology that has gained momentum over the last decade and generally describes network architectures that control the network using a combination of programmable switches with control logic implemented as a separate software program.

### 1.3.6 Internetworks

Many networks exist in the world, and they often use different hardware and software technologies. People connected to one network often want to communicate with people attached to a different one. The fulfillment of this desire requires that different, and frequently incompatible, networks be connected. A collection of interconnected networks is called an **internetwork** or **internet**. We will use these terms in a generic sense, in contrast to the global **Internet** (which is one specific internet), which we will always capitalize. The Internet connects content providers, access networks, enterprise networks, home networks, and many other networks to one another. We will look at the Internet in great detail later in this book.

A network comprises the combination of a subnet and its hosts. However, the word “network” is often used in a loose (and confusing) sense as well. A subnet might be described as a network, as in the case of the “ISP network” of Fig. 1-11.



An internetwork might also be described as a network, as in the case of the WAN in Fig. 1-9. We will follow similar practice, and if we are distinguishing a network from other arrangements, we will stick with our original definition of a collection of computers interconnected by a single technology.

An internet entails the interconnection of distinct, independently operated networks. In our view, connecting a LAN and a WAN or connecting two LANs is the usual way to form an internetwork, but there is little agreement over terminology in this area. Generally speaking, if two or more independently operated networks pay to interconnect, or if two or more networks use fundamentally different underlying technology (e.g., broadcast versus point-to-point and wired versus wireless), we probably have an internetwork.

The device that makes a connection between two or more networks and provides the necessary translation, both in terms of hardware and software, is a **gateway**. Gateways are distinguished by the layer at which they operate in the protocol hierarchy. We will have much more to say about layers and protocol hierarchies in the next section, but for now imagine that higher layers are more tied to applications, such as the Web, and lower layers are more tied to transmission links, such as Ethernet. Because the benefit of forming an internet is to connect computers across networks, we do not want to use too low-level a gateway or we will be unable to make connections between different kinds of networks. We do not want to use too high-level a gateway either, or the connection will only work for particular applications. The level in the middle that is “just right” is often called the network layer, and a router is a gateway that switches packets at the network layer. Generally speaking, an internetwork will be connected by network-layer gateways, or routers; however, even a single large network often contains many routers.

## 1.4 EXAMPLES OF NETWORKS

The subject of computer networking covers many different kinds of networks, large and small, well known and less well known. They have different goals, scales, and technologies. In the following sections, we will look at some examples, to get an idea of the variety one finds in the area of computer networking.

We will start with the Internet, probably the best-known “network,” and look at its history, evolution, and technology. Then, we will consider the mobile phone network. Technically, it is quite different from the Internet. Next, we will introduce IEEE 802.11, the dominant standard for wireless LANs.

### 1.4.1 The Internet

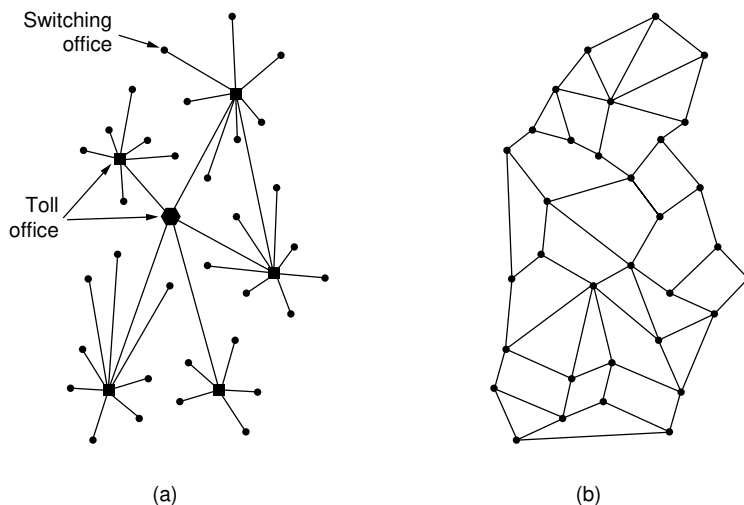
The Internet is a vast collection of different networks that use certain common protocols and provide certain common services. It is an unusual system in that it was not planned by any single organization, and it is not controlled by any single

organization, either. To better understand it, let us start from the beginning and see how it has developed and why. For a wonderful history of how the Internet developed, John Naughton's (2000) book is highly recommended. It is one of those rare books that is not only fun to read but also has 20 pages of *ibid.*'s and *op. cit.*'s for the serious historian. Some of the material in this section is based on this book. For a more recent history, try Brian McCullough's book (2018).

Of course, countless technical books have been written about the Internet, its history, and its protocols as well. For more information, see, for example, Severance (2015).

### The ARPANET

The story begins in the late 1950s. At the height of the Cold War, the U.S. DoD (Department of Defense) wanted a command-and-control network that could survive a nuclear war. At that time, all military communications used the public telephone network, which was considered vulnerable. The reason for this belief can be gleaned from Fig. 1-12(a). Here the black dots represent telephone switching offices, each of which was connected to thousands of telephones. These switching offices were, in turn, connected to higher-level switching offices (toll offices), to form a national hierarchy with only a small amount of redundancy. The vulnerability of the system was that the destruction of a few key toll offices could fragment it into many isolated islands so that generals in the Pentagon could not call a base in Los Angeles.



**Figure 1-12.** (a) Structure of the telephone system. (b) Baran's proposal.

Around 1960, the DoD awarded a contract to the RAND Corporation to find a solution. One of its employees, Paul Baran, came up with the highly distributed

and fault-tolerant design of Fig. 1-12(b). Since the paths between any two switching offices were now much longer than analog signals could travel without distortion, Baran proposed using digital packet-switching technology. Baran wrote several reports for the DoD describing his ideas in detail (Baran, 1964). Officials at the Pentagon liked the concept and asked AT&T, then the U.S.' national telephone monopoly, to build a prototype. AT&T dismissed Baran's ideas out of hand. The biggest and richest corporation in the world was not about to allow some young whippersnapper (out in California, no less—AT&T was then an East Coast company) tell it how to build a telephone system. They said Baran's network could not be built and the idea was killed.

Several years went by and still the DoD did not have a better command-and-control system. To understand what happened next, we have to go back all the way to October 1957, when the Soviet Union beat the U.S. into space with the launch of the first artificial satellite, Sputnik. When President Dwight Eisenhower tried to find out who was asleep at the switch, he was appalled to find the Army, Navy, and Air Force squabbling over the Pentagon's research budget. His immediate response was to create a single defense research organization, **ARPA**, the **Advanced Research Projects Agency**. ARPA had no scientists or laboratories; in fact, it had nothing more than an office and a small (by Pentagon standards) budget. It did its work by issuing grants and contracts to universities and companies whose ideas looked promising to it.

For the first few years, ARPA tried to figure out what its mission should be. In 1967, the attention of Larry Roberts, a program manager at ARPA who was trying to figure out how to provide remote access to computers, turned to networking. He contacted various experts to decide what to do. One of them, Wesley Clark, suggested building a packet-switched subnet, connecting each host to its own router.

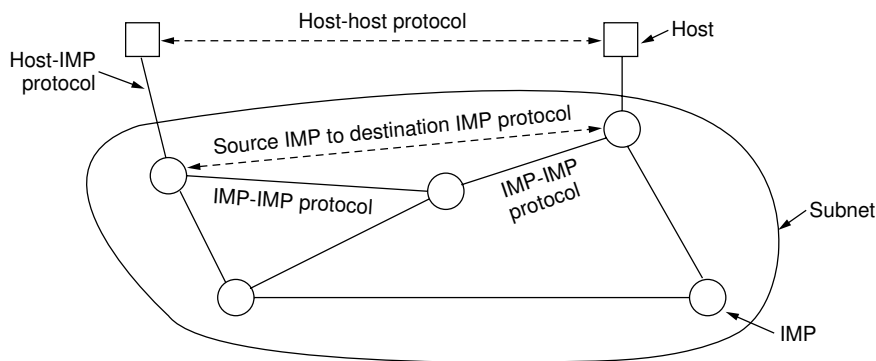
After some initial skepticism, Roberts bought the idea and presented a somewhat vague paper about it at the ACM SIGOPS Symposium on Operating System Principles held in Gatlinburg, Tennessee, in late 1967 (Roberts, 1967). Much to Roberts' surprise, another paper at the conference described a similar system that had not only been designed but actually fully implemented under the direction of Donald Davies at the National Physical Laboratory in England. The NPL system was not a national system by any means. It just connected several computers on the NPL campus. Nevertheless, it convinced Roberts that packet switching could be made to work. Furthermore, it cited Baran's now discarded earlier work. Roberts came away from Gatlinburg determined to build what later became known as the **ARPANET**.

In the plan that was developed, the subnet would consist of minicomputers called **IMPs (Interface Message Processors)** connected by then-state-of-the-art 56-kbps transmission lines. For high reliability, each IMP would be connected to at least two other IMPs. Each packet sent across the subnet was to contain the full destination address, so if some lines and IMPs were destroyed, subsequent packets could be automatically rerouted along alternative paths.

Each node of the network was to consist of an IMP and a host, in the same room, connected by a short wire. A host could send messages of up to 8063 bits to its IMP, which would then break these up into packets of at most 1008 bits and forward them independently toward the destination. Each packet was received in its entirety before being forwarded, so the subnet was the first electronic store-and-forward packet-switching network.

ARPA then put out a tender for building the subnet. Twelve companies bid for it. After evaluating all the proposals, ARPA selected BBN, a consulting firm based in Cambridge, Massachusetts, and in December 1968 awarded it a contract to build the subnet and write the subnet software. BBN chose to use specially modified Honeywell DDP-316 minicomputers with 12K 16-bit words of magnetic core memory as the IMPs. The IMPs did not have disks since moving parts were considered unreliable. The IMPs were interconnected by 56-kbps lines leased from telephone companies. Although 56 kbps is now often the only choice of people in rural areas, back then, it was the best money could buy.

The software was split into two parts: subnet and host. The subnet software consisted of the IMP end of the host-IMP connection, the IMP-IMP protocol, and a source IMP to destination IMP protocol designed to improve reliability. The original ARPANET design is shown in Fig. 1-13.



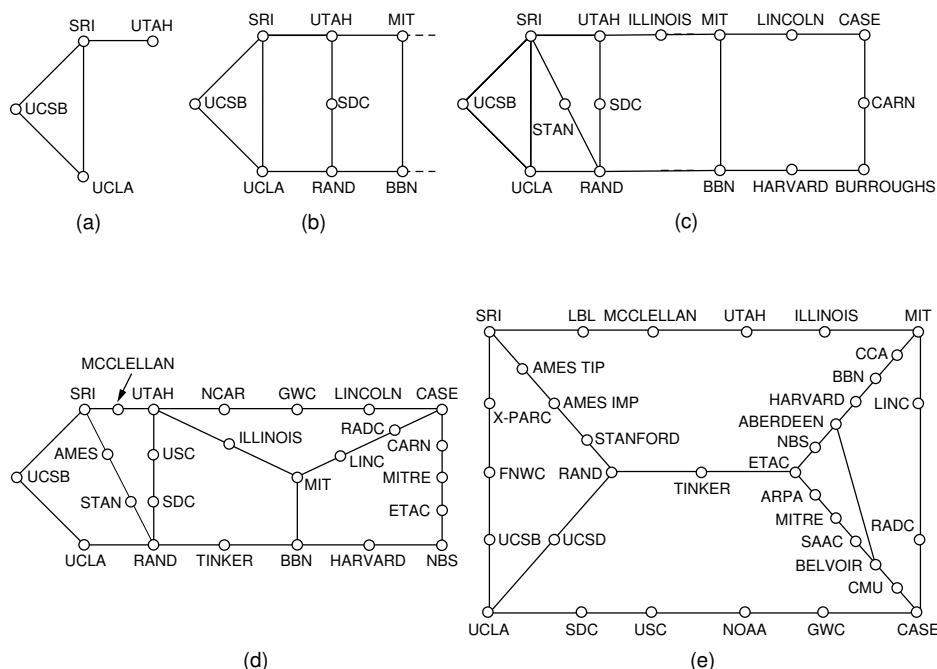
**Figure 1-13.** The original ARPANET design.

Outside the subnet, software was also needed, namely, the host end of the host-IMP connection, the host-host protocol, and the application software. It soon became clear that BBN was of the opinion that when it had accepted a message on a host-IMP wire and placed it on the host-IMP wire at the destination, its job was done.

Roberts had a problem, though: the hosts needed software too. To deal with it, he convened a meeting of network researchers, mostly graduate students, at Snowbird, Utah, in the summer of 1969. The graduate students expected some network

expert to explain the grand design of the network and its software to them and then assign each of them the job of writing part of it. They were astounded when there was no network expert and no grand design. They had to figure out what to do on their own.

Nevertheless, somehow an experimental network went online in December 1969 with four nodes: at UCLA, UCSB, SRI, and the University of Utah. These four were chosen because all had a large number of ARPA contracts, and all had different and completely incompatible host computers (just to make it more fun). The first host-to-host message had been sent two months earlier from the UCLA node by a team led by Len Kleinrock (a pioneer of the theory of packet switching) to the SRI node. The network grew quickly as more IMPs were delivered and installed; it soon spanned the United States. Figure 1-14 shows how rapidly the ARPANET grew in the first 3 years.



**Figure 1-14.** Growth of the ARPANET. (a) December 1969. (b) July 1970. (c) March 1971. (d) April 1972. (e) September 1972.

In addition to helping the fledgling ARPANET grow, ARPA also funded research on the use of satellite networks and mobile packet radio networks. In one now-famous demonstration, a big truck driving around in California used the packet radio network to send messages to SRI, which were then forwarded over the ARPANET to the East Coast, where they were then shipped to University College

in London over the satellite network. This allowed a researcher in the truck to use a computer in London while driving around in California.

This experiment also demonstrated that the existing ARPANET protocols were not suitable for running over different networks. This observation led to more research on protocols, culminating with the invention of the TCP/IP protocols (Cerf and Kahn, 1974). TCP/IP was specifically designed to handle communication over internetworks, something becoming increasingly important as more and more networks were hooked up to the ARPANET.

To encourage adoption of these new protocols, ARPA awarded several contracts to implement TCP/IP on different computer platforms, including IBM, DEC, and HP systems, as well as for Berkeley UNIX. Researchers at the University of California at Berkeley rewrote TCP/IP with a new programming interface called **sockets** for the upcoming 4.2BSD release of Berkeley UNIX. They also wrote many application, utility, and management programs to show how convenient it was to use the network with sockets.

The timing was perfect. Many universities had just acquired a second or third VAX computer and a LAN to connect them, but they had no networking software. When 4.2BSD came along, with TCP/IP, sockets, and many network utilities, the complete package was adopted immediately. Furthermore, with TCP/IP, it was easy for the LANs to connect to the ARPANET, and many did. As a result, TCP/IP use grew rapidly during the mid-1970s.

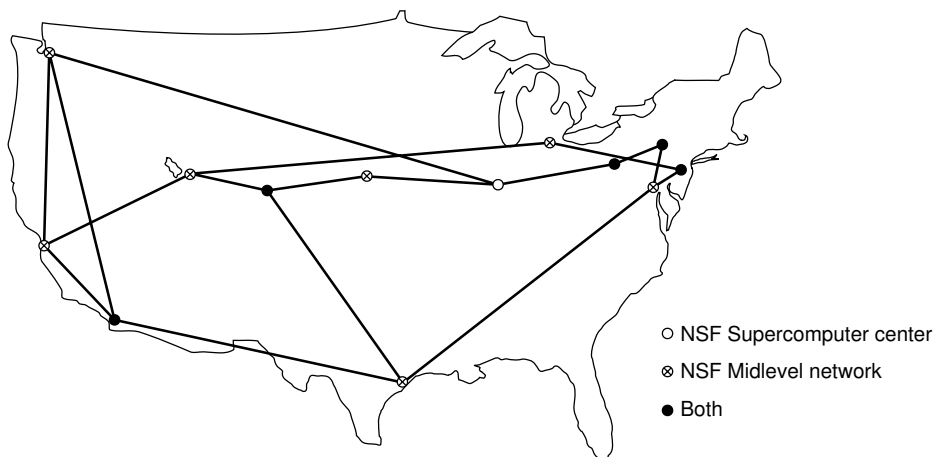
## NSFNET

By the late 1970s, NSF (the U.S. National Science Foundation) saw the enormous impact the ARPANET was having on university research, allowing scientists across the country to share data and collaborate on research projects. However, to get on the ARPANET a university had to have a research contract with the DoD. Many did not have a contract. NSF's initial response was to fund **CSNET (Computer Science Network)** in 1981. It connected computer science departments and industrial research labs to the ARPANET via dial-up and leased lines. In the late 1980s, the NSF went further and decided to design a successor to the ARPANET that would be open to all university research groups.

To have something concrete to start with, NSF decided to build a backbone network to connect its six supercomputer centers, in San Diego, Boulder, Champaign, Pittsburgh, Ithaca, and Princeton. Each supercomputer was given a little brother, consisting of an LSI-11 microcomputer called a **fuzzball**. The fuzzballs were connected with 56-kbps leased lines and formed the subnet, the same hardware technology the ARPANET used. The software technology was different, however: the fuzzballs spoke TCP/IP right from the start, making it the first TCP/IP WAN.

NSF also funded some (eventually about 20) regional networks that connected to the backbone to allow users at thousands of universities, research labs, libraries,

and museums to access any of the supercomputers and to communicate with one another. The complete network, including backbone and the regional networks, was called **NSFNET (National Science Foundation Network)**. It connected to the ARPANET through a link between an IMP and a fuzzball in the Carnegie-Mellon machine room. The first NSFNET backbone is illustrated in Fig. 1-15 superimposed on a map of the United States.



**Figure 1-15.** The NSFNET backbone in 1988.

NSFNET was an instantaneous success and was overloaded from the word go. NSF immediately began planning its successor and awarded a contract to the Michigan-based MERIT consortium to run it. Fiber optic channels at 448 kbps were leased from MCI (which was purchased by Verizon in 2006) to provide the version 2 backbone. IBM PC-RTs were used as routers. This, too, was soon overwhelmed, and by 1990, the second backbone was upgraded to 1.5 Mbps.

As growth continued, NSF realized that the government could not continue financing networking forever. Furthermore, commercial organizations wanted to join but were forbidden by NSF's charter from using networks NSF paid for. Consequently, NSF encouraged MERIT, MCI, and IBM to form a nonprofit corporation, **ANS (Advanced Networks and Services)**, as the first step along the road to commercialization. In 1990, ANS took over NSFNET and upgraded the 1.5-Mbps links to 45 Mbps to form **ANSNET**. This network operated for 5 years and was then sold to America Online. But by then, various companies were offering commercial IP service and it was clear that the government should now get out of the networking business.

To ease the transition and make sure every regional network could communicate with every other regional network, NSF awarded contracts to four different network operators to establish a **NAP (Network Access Point)**. These operators